

Profiling the Memory Usage of OCaml Applications without Changing their Behavior

OCaml 2013

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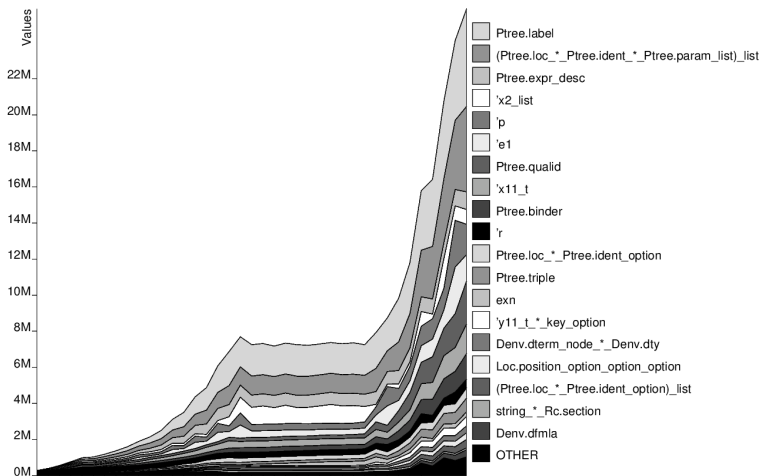
Sept. 24 2013 - Boston (MA)

Memory Problems

- What ?
 - Study the memory behavior of OCaml programs
 - Memory profiling tools

- Why ?
 - To decrease memory footprint
 - To fix memory leaks
 - To spend less time in memory management

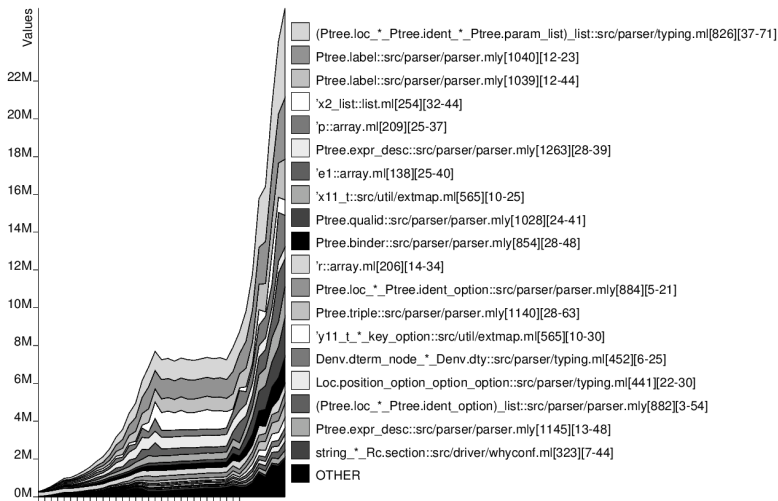
Real World Example – Why3¹ (1/2)



¹Why3 is a platform for deductive program verification
(<http://why3.lri.fr/>)

Real World Example – Why3 (2/2)

With locations precision



How do we do that ?

```
$ opam switch 4.00.1+ocp-memprof
```

```
$ opam install why3
```

```
$ OCAMLRUNPARAM=m why3replayer.opt -C why3.conf p9_16
```

this step will generate a lot of snapshots of heap image

No need to change your code nor the compilation options.

No impact on execution time.

```
$ opam install ocp-memprof
```

```
$ ocp-memprof -loc -sizes PID
```

this step analyzes all these snapshots

Look at the graphs.

Snapshots

What is a snapshot ?

- Compressed version of the heap
- Location identifiers, graph with pointers, etc.
- Save globals (toplevel modules)

How do we obtain these snapshots ?

- Computed by a linear scan of all chunks² which contain sets of consecutive blocks.

² huge block of memory

Generate A Snapshot

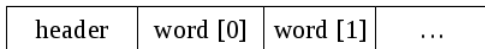
Two ways to trigger the generation of snapshots

- Use `OCAMLRUNPARAM=m` force a program to generate a snapshot after every GC
- Request explicitly the program to generate a snapshot
 - by sending a HUP signal (very useful for server-like application, cf `mldonkey`)
 - in module `GC`, use the following function

```
val dump_heap : string -> unit
```

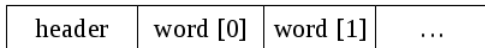
Patched Compiler (1/3)

OCaml memory block:

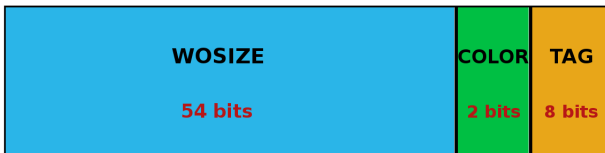


Patched Compiler (1/3)

OCaml memory block:

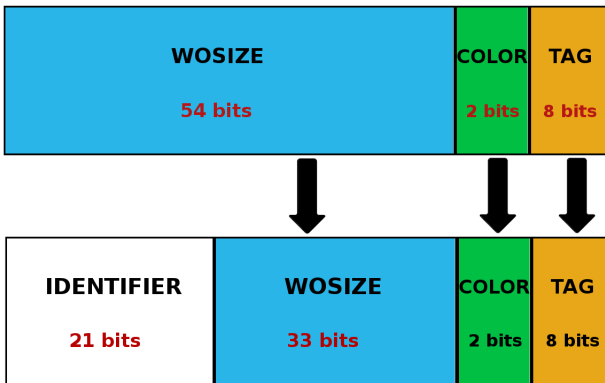


OCaml block's header (one word) on 64-bit machines:



Patched Compiler (2/3)

Header after our modification:



Patched Compiler (3/3)



- Minimal impact on performance (only when generating snapshots)

Patched Compiler (3/3)



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- No space overhead

Patched Compiler (3/3)



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- No impact on GC (its behavior is not changed)

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- Only on 64-bit platforms
- Location identifiers are limited ($2^{21} \sim 2$ million allocation sites)

Patched Compiler (3/3)

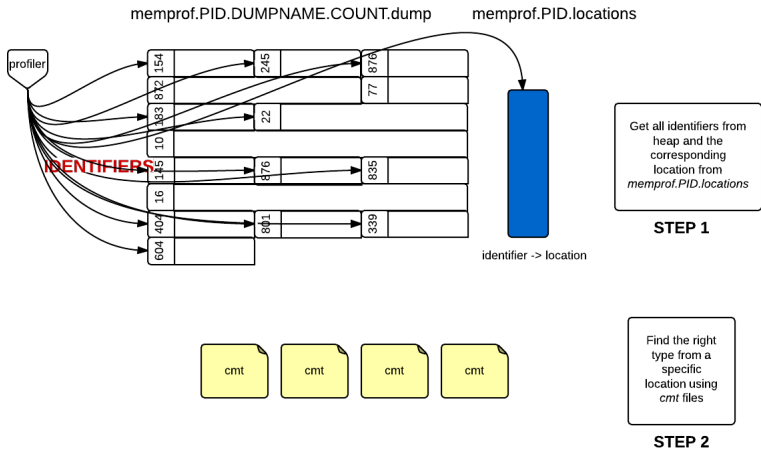


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- Only on 64-bit platforms
- Location identifiers are limited ($2^{21} \sim 2$ million allocation sites)
- Maximum block size is now 64GB

One Tool Based On Identifiers



*A cmt file is a binary file containing the typed AST

Conclusion

Future Work:

- Improve the current framework
 - Aggregate information by type and location (work in progress)
 - Recover more types (e.g. using G.Henry's work)
 - Display life span of values (number of GC for example)
- More tools based to analyzed snapshots:
 - a graphical assistant to explore snapshots
 - a tool which use pointers to see which root retains some specific values

Questions ?